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# Regret Analysis of the Anytime Optimally Confident UCB Algorithm

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**Tor Lattimore**

Department of Computing Science  
University of Alberta, Canada  
tor.lattimore@gmail.com

## Abstract

I introduce and analyse an anytime version of the Optimally Confident UCB (OCUCB) algorithm designed for minimising the cumulative regret in finite-armed stochastic bandits with subgaussian noise. The new algorithm is simple, intuitive (in hindsight) and comes with the strongest finite-time regret guarantees for a horizon-free algorithm so far. I also show a finite-time lower bound that nearly matches the upper bound and has a trivial proof.

## 1 Introduction

A potential drawback of the Optimally Confident UCB algorithm for finite-armed bandits is that it depends on advance knowledge of the horizon [Lattimore, 2015]. I address this issue by analysing an anytime version that performs about as well empirically and for which the regret guarantee suffers only an additional  $\log \log(n)$  additive term. The proof is also significantly easier – though still quite involved – and more importantly is quite intuitive. For the sake of brevity I will give neither a detailed introduction nor an exhaustive survey of the literature. Readers looking for a gentle primer on multi-armed bandits might enjoy the monograph by [Bubeck and Cesa-Bianchi \[2012\]](#) from which I borrow notation. Let  $K$  be the number of arms and  $I_t \in \{1, \dots, K\}$  be the arm chosen in round  $t$ . The reward is  $X_t = \mu_{I_t} + \xi_t$  where  $\mu \in \mathbb{R}^K$  is the unknown vector of means and the noise term  $\xi_t$  is assumed to be 1-subgaussian (therefore zero-mean). The  $n$ -step pseudo-regret of strategy  $\pi$  given mean vector  $\mu$  with maximum mean  $\mu^* = \max_i \mu_i$  is

$$R_\mu^\pi(n) = n\mu^* - \mathbb{E} \sum_{t=1}^n \mu_{I_t},$$

where the expectation is taken with respect to the uncertainty in both the rewards and actions. In all analysis I make the standard notational assumption that  $\mu_1 \geq \mu_2 \geq \dots \geq \mu_K$ . The new algorithm is called OCUCB- $n$  and depends on two parameters  $\eta > 1$  and  $\rho \in [1/2, 1]$ . The algorithm chooses  $I_t = t$  in rounds  $t \leq K$  and subsequently  $I_t = \arg \max_i \gamma_i(t)$  with

$$\gamma_i(t) = \hat{\mu}_i(t-1) + \sqrt{\frac{2\eta \log(B_i(t-1))}{T_i(t-1)}}, \quad (1)$$

where  $T_i(t-1)$  is the number of times arm  $i$  has been chosen after round  $t-1$  and  $\hat{\mu}_i(t-1)$  is the empirical estimate of its mean and

$$B_i(t-1) = \max \left\{ e, \log(t), t \log(t) \left( \sum_{j=1}^K \min \{T_j(t-1), T_j(t-1)^\rho T_i(t-1)^{1-\rho}\} \right)^{-1} \right\}.$$

Besides the algorithm, the contribution of this article is a proof that OCUCB- $n$  satisfies a nearly optimal regret bound.

**Theorem 1.** *If  $\rho \in [1/2, 1]$  and  $\eta > 1$ , then*

$$R_{\mu}^{\text{OCUCB-}n}(n) \leq C_{\eta} \sum_{i:\Delta_i>0} \left( \Delta_i + \frac{1}{\Delta_i} \log \max \left\{ \frac{n\Delta_i^2 \log(n)}{k_{i,\rho}}, \log(n) \right\} \right),$$

where  $\Delta_i = \mu^* - \mu_i$  and  $k_{i,\rho} = \sum_{j=1}^K \min\{1, \Delta_i^{2\rho}/\Delta_j^{2\rho}\}$  and  $C_{\eta} > 0$  is a constant that depends only on  $\eta$ . Furthermore, for all  $\rho \in [0, 1]$  it holds that  $\limsup_{n \rightarrow \infty} R_{\mu}^{\text{OCUCB-}n}(n)/\log(n) \leq \sum_{i:\Delta_i>0} \frac{2\eta}{\Delta_i}$ .

Asymptotically the upper bound matches the lower bound given by [Lai and Robbins \[1985\]](#) except for a factor of  $\eta$ . In the non-asymptotic regime the additional terms inside the logarithm significantly improves on UCB. The bound in [Theorem 1](#) corresponds to a worst-case regret that is suboptimal by a factor of just  $\sqrt{\log \log n}$ . Algorithms achieving the minimax rate are [MOSS \[Audibert and Bubeck, 2009\]](#) and [OCUCB](#), but both require advance knowledge of the horizon. The quantity  $k_{i,\rho} \in [1, K]$  may be interpreted as the number of ‘‘effective’’ arms with larger values leading to improved regret. A simple observation is that  $k_{i,\rho}$  is always non-increasing in  $\rho$ , which makes  $\rho = 1/2$  the canonical choice. In the special case that all suboptimal arms have the same expected payoff, then  $k_{i,\rho} = K$  for all  $\rho$ . Interestingly I could not find a regime for which the algorithm is empirically sensitive to  $\rho \in [1/2, 1]$ . If  $\rho = 1$ , then except for  $\log \log$  additive terms the problem dependent regret enjoyed by [OCUCB- \$n\$](#)  is equivalent to [OCUCB](#). Finally, if  $\rho = 0$ , then the asymptotic result above applies, but the algorithm in that case essentially reduces to [MOSS](#), which is known to suffer suboptimal finite-time regret in certain regimes [\[Lattimore, 2015\]](#).

**Intuition for regret bound.** Let us fix a strategy  $\pi$  and mean vector  $\mu \in \mathbb{R}^K$  and suboptimal arm  $i$ . Suppose that  $\mathbb{E}[T_i(n)] \leq \Delta_i^{-2} \log(1/\delta)/2$  for some  $\delta \in (0, 1)$ . Now consider the alternative mean reward  $\mu'$  with  $\mu'_j = \mu_j$  for  $j \neq i$  and  $\mu'_i = \mu_i + 2\Delta_i$ , which means that  $i$  is the optimal action for mean vector  $\mu'$ . Standard information-theoretic analysis shows that  $\mu$  and  $\mu'$  are not statistically separable at confidence level  $\delta$  and in particular, if  $\Delta_i$  is large enough, then  $R_{\mu'}^{\pi}(n) = \Omega(n\delta\Delta_i)$ . For mean  $\mu'$  we have  $\Delta'_j = \mu'_i - \mu'_j \approx \max\{\Delta_i, \Delta_j\}$  and for any reasonable algorithm we would like

$$\sum_{j:\Delta'_j>0} \frac{\log(n)}{\Delta'_j} \gtrsim R_{\mu'}^{\pi}(n) = \Omega(n\delta\Delta_i).$$

But this implies that  $\delta$  should be chosen such that

$$\delta = O\left(\frac{\log(n)}{n} \sum_{j:\Delta'_j>0} \frac{1}{\Delta'_j \Delta_i}\right) = O\left(\frac{\log(n)k_{i,1/2}}{n\Delta_i^2}\right),$$

which up to  $\log \log$  terms justifies the near-optimality of the regret guarantee given in [Theorem 1](#) for  $\rho$  close to  $1/2$ . Of course  $\Delta$  is not known in advance, so no algorithm can choose this confidence level. The trick is to notice that arms  $j$  with  $\Delta_j \leq \Delta_i$  should be played about as often as arm  $i$  and arms  $j$  with  $\Delta_j > \Delta_i$  should be played about as much as arm  $i$  until  $T_j(t-1) \approx \Delta_j^{-2}$ . This means that as  $T_i(t-1)$  approaches the critical number of samples  $\Delta_i^{-2}$  we can approximate

$$\sum_{j=1}^K \min \left\{ T_i(t-1), T_j(t-1)^{\frac{1}{2}} T_i(t-1)^{\frac{1}{2}} \right\} \approx \sum_{j=1}^K \min \left\{ \Delta_i^{-2}, \Delta_j^{-1} \Delta_i^{-1} \right\} = \frac{k_{i,1/2}}{\Delta_i^2}.$$

Then the index used by [OCUCB- \$n\$](#)  is justified by ignoring  $\log \log$  terms and the usual  $n \approx t$  used by UCB and other algorithms. [Theorem 1](#) is proven by making the above approximation rigorous. The argument for this choice of confidence level is made concrete in [Appendix A](#) where I present a lower bound that matches the upper bound except for  $\log \log(n)$  additive terms.

## 2 Concentration

The regret guarantees rely on a number of concentration inequalities. For this section only let  $X_1, X_2, \dots$  be i.i.d. 1-subgaussian and  $S_n = \sum_{t=1}^n X_t$  and  $\hat{\mu}_n = S_n/n$ . The first lemma below is well known and follows trivially from the maximal inequality.

**Important remark.** For brevity I use  $O_\eta(1)$  to indicate a constant that depends on  $\eta$  but not other variables such as  $n$  and  $\mu$ . The dependence is never worse than polynomial in  $1/(\eta - 1)$ .

**Lemma 2.** If  $\varepsilon > 0$ , then  $\mathbb{P}\{\exists t \leq n : S_t \geq \varepsilon\} \leq \exp\left(-\frac{\varepsilon^2}{2n}\right)$ .

The following lemma analyses the likelihood that  $S_n$  ever exceeds  $f(n) = \sqrt{2\eta n \log \max\{e, \log n\}}$  where  $\eta > 1$ . By the law of the iterated logarithm  $\limsup_{n \rightarrow \infty} S_n/f(n) = \sqrt{1/\eta}$  a.s. and for small  $\delta$  it has been shown by [Garivier \[2013\]](#) that

$$\mathbb{P}\left\{\exists n : S_n \geq \sqrt{2n \log\left(\frac{\log(n)}{\delta}\right)}\right\} = O(\delta).$$

The case where  $\delta = \Omega(1)$  seems not to have been analysed and relies on the usual peeling trick, but without the union bound.

**Lemma 3.** There exists a monotone non-decreasing function  $p : (1, \infty) \rightarrow (0, 1]$  such that for all  $\eta > 1$  it holds that  $\mathbb{P}\left\{\forall n : S_n \leq \sqrt{2\eta n \log \max\{e, \log n\}}\right\} \geq p(\eta)$ .

**Lemma 4.** Let  $b \geq 1$  and  $\Delta > 0$  and  $\tau = \min\left\{n : \sup_{t \geq n} \hat{\mu}_t + \sqrt{\frac{2\eta \log(b)}{t}} < \Delta\right\}$ , then

$$\mathbb{E}[\tau] \leq \sqrt{\mathbb{E}[\tau^2]} = O_\eta(1) \cdot \left(1 + \frac{1}{\Delta^2} \log_+(b)\right) \quad \text{where } \log_+(x) = \max\{1, \log(x)\}.$$

The final concentration lemma is quite powerful and forms the lynch-pin of the following analysis.

**Lemma 5.** Let  $\Delta > 0$  and  $\rho \in [0, 1]$  and  $d \in \{1, 2, \dots\}$  and  $\lambda_1, \dots, \lambda_d \in [1, \infty]$  be constants. Furthermore, let  $\alpha$  be the random variable given by

$$\alpha = \inf\left\{\alpha \geq 0 : \inf_s \hat{\mu}_s + \sqrt{\frac{2\eta}{s} \log \max\left\{1, \frac{\alpha}{\sum_{i=1}^d \min\{s, \lambda_i^\rho s^{1-\rho}\}}\right\}} \geq -\Delta\right\}.$$

Finally let  $\beta = \inf\{\beta \geq 0 : \beta \log(\beta) = \alpha\}$ . Then

$$(a) \text{ If } \rho \in (1/2, 1], \text{ then } \Delta \mathbb{E}[\alpha] = O\left(\frac{1}{(2\rho - 1)(\eta - 1)^2}\right) \cdot \sum_{i=1}^d \min\{\Delta^{-1}, \sqrt{\lambda_i}\}$$

$$(b) \text{ If } \rho \in [1/2, 1], \text{ then } \Delta \mathbb{E}[\beta] = O\left(\frac{1}{(\eta - 1)^2}\right) \cdot \sum_{i=1}^d \min\{\Delta^{-1}, \sqrt{\lambda_i}\}$$

The proofs of Lemmas 3 to 5 may be found in Appendices B to D.

### 3 Analysis of the KL-UCB+ Algorithm

Let us warm up by analysing a simpler algorithm, which chooses the arm that maximises the following index.

$$\gamma_i(t) = \hat{\mu}_i(t-1) + \sqrt{\frac{2\eta}{T_i(t-1)} \log\left(\frac{t}{T_i(t-1)}\right)}. \quad (2)$$

Strategies similar to this have been called KL-UCB+ and suggested as a heuristic by [Garivier and Cappé \[2011\]](#) (this version is specified to the subgaussian noise model). Recently [Kaufmann \[2016\]](#) has established the asymptotic optimality of strategies with approximately this form, but finite-time analysis has not been available until now. Bounding the regret will follow the standard path of bounding  $\mathbb{E}[T_i(n)]$  for each suboptimal arm  $i$ . Let  $\hat{\mu}_{i,s}$  be the empirical estimate of the mean of the

$i$ th arm having observed  $s$  samples. Define  $\tau_i$  and  $\tau_\Delta$  by

$$\tau_i = \min \left\{ t \geq 1/\Delta_i^2 : \sup_{s \geq t} \hat{\mu}_{i,s} + \sqrt{\frac{2\eta}{s} \log \max\{1, n\Delta_i^2\}} < \mu_i + \frac{\Delta_i}{2} \right\}$$

$$\tau_\Delta = \min \left\{ t : \inf_{1 \leq s \leq n} \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \max\left\{1, \frac{t}{s}\right\}} \geq \mu_1 - \frac{\Delta_i}{2} \right\}.$$

If  $T_i(t-1) \geq \tau_i$  and  $t \geq \tau_\Delta$ , then by the definition of  $\tau_\Delta$  we have  $\gamma_1(t) \geq \mu_i + \Delta_i/2$  and by the definition of  $\tau_i$

$$\gamma_i(t) = \hat{\mu}_i(t-1) + \sqrt{\frac{2\eta \log(t/T_i(t-1))}{T_i(t-1)}} \leq \hat{\mu}_i(t-1) + \sqrt{\frac{2\eta \log(n\Delta_i^2)}{T_i(t-1)}} < \mu_i + \frac{\Delta_i}{2},$$

which means that  $I_t \neq i$ . Therefore  $T_i(n)$  may be bounded in terms of  $\tau_i$  and  $\tau_\Delta$  as follows:

$$T_i(n) = \sum_{t=1}^n \mathbb{1}\{I_t = i\} \leq \tau_\Delta + \sum_{t=\tau_\Delta+1}^n \mathbb{1}\{I_t = i \text{ and } T_i(t-1) < \tau_i\} \leq \tau_i + \tau_\Delta.$$

It remains to bound the expectations of  $\tau_i$  and  $\tau_\Delta$ . By Lemma 5a with  $d = 1$  and  $\rho = 1$  and  $\lambda_1 = \infty$  it follows that  $\mathbb{E}[\tau_\Delta] = O_\eta(1) \cdot \Delta_i^{-2}$  and by Lemma 4

$$\mathbb{E}[\tau_i] = O_\eta(1) \cdot \left( 1 + \frac{1}{\Delta_i^2} \log_+(n\Delta_i^2) \right).$$

Therefore the strategy in Eq. (2) satisfies:

$$R_\mu^{\text{KL-UCB}^+}(n) = \sum_{i:\Delta_i>0} \Delta_i \mathbb{E}[T_i(n)] = O_\eta(1) \cdot \sum_{i:\Delta_i>0} \left( \Delta_i + \frac{1}{\Delta_i} \log_+(n\Delta_i^2) \right).$$

**Remark 6.** Without changing the algorithm and by optimising the constants in the proof it is possible to show that  $\limsup_{n \rightarrow \infty} R_\mu^{\text{KL-UCB}^+}(n)/\log(n) \leq \sum_{i:\Delta_i>0} 2\eta/\Delta_i$ , which is just a factor of  $\eta$  away from the asymptotic lower bound of [Lai and Robbins \[1985\]](#).

## 4 Proof of Theorem 1

The proof follows along similar lines as the warm-up, but each step becomes more challenging, especially controlling  $\tau_\Delta$ .

### Step 1: Setup and preliminary lemmas

Define  $\Phi$  to be the random set of arms for which the empirical estimate never drops below the critical boundary given by the law of iterated logarithm.

$$\Phi = \left\{ i > 2 : \hat{\mu}_{i,s} + \sqrt{\frac{2\eta_1 \log \max\{e, \log s\}}{s}} \geq \mu_i \text{ for all } s \right\}, \quad (3)$$

where  $\eta_1 = (1 + \eta)/2 \in (1, \eta)$ . By Lemma 3,  $\mathbb{P}\{i \in \Phi\} \geq p(\eta_1) > 0$ . It will be important that  $\Phi$  only includes arms  $i > 2$  and that the events  $i, j \in \Phi$  are independent for  $i \neq j$ . From the definition of the index  $\gamma$  and for  $i \in \Phi$  it holds that  $\gamma_i(t) > \mu_i$  for all  $t$ . The following lemma shows that the pull counts for optimistic arms “chase” those of other arms up the point that they become clearly suboptimal.

**Lemma 7.** *There exists a constant  $c_\eta \in (0, 1)$  depending only on  $\eta$  such that if (a)  $j \in \Phi$  and (b)  $\hat{\mu}_i(t-1) \leq \mu_i + \Delta_i/2$  and (c)  $T_j(t-1) \leq c_\eta \min\{T_i(t-1), \Delta_j^{-2}\}$ , then  $I_t \neq i$ .*

*Proof.* First note that  $T_j(t-1) \leq T_i(t-1)$  implies that  $B_j(t-1) \geq B_i(t-1)$ . Comparing the indices:

$$\gamma_i(t) = \hat{\mu}_i(t-1) + \sqrt{\frac{2\eta \log B_i(t-1)}{T_i(t-1)}} \leq \mu_i + \sqrt{\frac{2\eta c_\eta \log B_j(t-1)}{T_j(t-1)}} + \frac{\Delta_i}{2}.$$

On the other hand, by choosing  $c_\eta$  small enough and by the definition of  $j \in \Phi$ :

$$\begin{aligned}\gamma_j(t) &= \hat{\mu}_j(t-1) + \sqrt{\frac{2\eta \log B_j(t-1)}{T_j(t-1)}} \geq \mu_j + \sqrt{\frac{2\eta c_\eta \log B_j(t-1)}{T_j(t-1)}} + \sqrt{\frac{c_\eta}{T_j(t-1)}} \\ &\geq \mu_1 + \sqrt{\frac{2\eta c_\eta \log B_j(t-1)}{T_j(t-1)}} > \gamma_i(t),\end{aligned}$$

which implies that  $I_t \neq i$ .  $\square$

Let  $J = \min \Phi$  be the optimistic arm with the largest return where if  $\Phi = \emptyset$  we define  $J = K + 1$  and  $\Delta_J = \max_i \Delta_i$ . By Lemma 3,  $i \in \Phi$  with constant probability, which means that  $J$  is sub-exponentially distributed with rate dependent on  $\eta$  only. Define  $K_{i,\rho}$  by

$$K_{i,\rho} = 1 + c_\eta \sum_{j \in \Phi, j \neq i} \min \left\{ 1, \frac{\Delta_i^{2\rho}}{\Delta_j^{2\rho}} \right\}, \quad (4)$$

where  $c_\eta$  is as chosen in Lemma 7. Since  $\mathbb{P}\{i \in \Phi\} = \Omega_\eta(1)$  we will have  $K_{i,\rho} = \Omega_\eta(k_{i,\rho})$  with high probability (this will be made formal later). Let

$$\begin{aligned}b_i &= \max \left\{ \frac{n\Delta_i^2 \log(n)}{k_{i,\rho}}, \log(n), e \right\} \quad \text{and} \quad B_i = \max \left\{ \frac{n\Delta_i^2 \log(n)}{K_{i,\rho}}, \log(n), e \right\} \\ \tau_i &= \min \left\{ s \geq \frac{1}{\Delta_i^2} : \sup_{s' \geq s} \hat{\mu}_{i,s'} + \sqrt{\frac{2\eta}{s'}} \log(B_i) \leq \mu_i + \frac{\Delta_i}{2} \right\}.\end{aligned} \quad (5)$$

The following lemma essentially follows from Lemma 4 and the fact that  $J$  is sub-exponentially distributed. Care must be taken because  $J$  and  $\tau_i$  are not independent. The proof is found in Appendix E.

**Lemma 8.**  $\mathbb{E}[\tau_i] \leq \mathbb{E}[J\tau_i] = O_\eta(1) \cdot \left( 1 + \frac{1}{\Delta_i^2} \log(b_i) \right)$ .

The last lemma in this section shows that if  $T_i(t-1) \geq \tau_i$ , then either  $i$  is not chosen or the index of the  $i$ th arm is not too large.

**Lemma 9.** *If  $T_i(t-1) \geq \tau_i$ , then  $I_t \neq i$  or  $\gamma_i(t) < \mu_i + \Delta_i/2$ .*

*Proof.* By the definition of  $\tau_i$  we have  $\tau_i \geq \Delta_i^{-2}$  and  $\hat{\mu}_i(t-1) \leq \mu_i + \Delta_i/2$ . By Lemma 7, if  $j \in \Phi$  and  $T_j(t-1) \leq c_\eta \min \{\Delta_i^{-2}, \Delta_j^{-2}\}$ , then  $I_t \neq i$ . Now suppose that  $T_j(t-1) \geq c_\eta \min \{\Delta_i^{-2}, \Delta_j^{-2}\}$  for all  $j \in \Phi$ . Then

$$\begin{aligned}B_i(t-1) &= \max \left\{ e, \log(t), t \log(t) \left( \sum_{j=1}^K \min \{T_i(t-1), T_j(t-1)^\rho T_i(t-1)^{1-\rho}\} \right)^{-1} \right\} \\ &\leq \max \left\{ e, \log(n), \frac{n\Delta_i^2 \log(n)}{K_{i,\rho}} \right\} = B_i.\end{aligned}$$

Therefore from the definition of  $\tau_i$  we have that  $\gamma_i(t) < \mu_i + \Delta_i/2$ .  $\square$

## Step 2: Regret decomposition

By Lemma 9, if  $T_i(n) \geq \tau_i$ , then  $I_t \neq i$  or  $\gamma_i(t) < \mu_i + \Delta_i/2$ . Now we must show there exists a  $j$  for which  $\gamma_j(t) \geq \mu_i + \Delta_i/2$ . This is true for arms  $i$  with  $\Delta_i \geq 2\Delta_J$  since by definition  $\gamma_J(t) > \mu_J \geq \mu_i + \Delta_i/2$  for all  $t$ . For the remaining arms we follow the idea used in Section 3 and define a random time for each  $\Delta > 0$ .

$$\tau_\Delta = \min \left\{ t : \inf_{s \geq t} \sup_j \gamma_j(s) \geq \mu_1 - \frac{\Delta}{2} \right\}. \quad (6)$$

Then the regret is decomposed as follows

$$R_\mu^{\text{OCUCB-}n}(n) \leq \mathbb{E} \left[ \sum_{i:\Delta_i > 0} \Delta_i \tau_i + 2\Delta_J \tau_{\Delta_J/4} + \sum_{i:\Delta_i < \Delta_J/4} \Delta_i \tau_{\Delta_i} \right]. \quad (7)$$

The next step is to show that the first sum is dominant in the above decomposition, which will lead to the result via Lemma 8 to bound  $\mathbb{E}[\Delta_i \tau_i]$ .

### Step 3: Bounding $\tau_\Delta$

This step is broken into two quite technical parts as summarised in the following lemma. The proofs of both results are quite similar, but the second is more intricate and is given in Appendix G.

**Lemma 10.** *The following hold:*

$$(a). \quad \mathbb{E} [\Delta_J \tau_{\Delta_J/4}] \leq O_\eta(1) \cdot \sum_{i:\Delta_i > 0} \sqrt{1 + \frac{\log(b_i)}{\Delta_i^2}}$$

$$(b). \quad \mathbb{E} \left[ \sum_{i:\Delta_i < \Delta_J/4} \Delta_i \tau_{\Delta_i} \right] \leq O_\eta(1) \cdot \sum_{i:\Delta_i > 0} \sqrt{1 + \frac{\log(b_i)}{\Delta_i^2}}.$$

*Proof of Lemma 10a.* Preparing to use Lemma 5, let  $\lambda \in (0, \infty]^K$  be given by  $\lambda_i = \tau_i$  for  $i$  with  $\Delta_i \geq 2\Delta_J$  and  $\lambda_i = \infty$  otherwise. Now define random variable  $\alpha$  by

$$\alpha = \inf \left\{ \alpha \geq 0 : \inf_s \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{\alpha}{\sum_{i=1}^K \min \{s, \lambda_i^\rho s^{1-\rho}\}} \right\}} \geq \mu_1 - \frac{\Delta_J}{8} \right\}$$

and  $\beta = \min \{ \beta \geq 0 : \beta \log(\beta) = \alpha \}$ . Then for  $t \geq \beta$  and abbreviating  $s = T_1(t-1)$  we have

$$\begin{aligned} \gamma_1(t) &= \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log B_1(t-1)} \\ &= \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \left( \max \left\{ e, \log(t), \frac{t \log(t)}{\sum_{i=1}^K \min \{s, T_i(t-1)^\rho s^{1-\rho}\}} \right\} \right)} \\ &\geq \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{\alpha}{\sum_{i=1}^K \min \{s, T_i(t-1)^\rho s^{1-\rho}\}} \right\}} \\ &\geq \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{\alpha}{\sum_{i=1}^K \min \{s, \lambda_i^\rho s^{1-\rho}\}} \right\}} \geq \mu_1 - \frac{\Delta_J}{8}, \end{aligned}$$

where the second last inequality follows since for arms with  $\Delta_i \geq 2\Delta_J$  we have  $T_i(n) \leq \tau_i = \lambda_i$  and for other arms  $\lambda_i = \infty$  by definition. The last inequality follows from the definition of  $\alpha$ . Therefore  $\tau_{\Delta_J/4} \leq \beta$  and so  $\mathbb{E}[\Delta_J \tau_{\Delta_J/4}] \leq \mathbb{E}[\Delta_J \beta]$ , which by Lemma 5b is bounded by

$$\begin{aligned} \mathbb{E}[\Delta_J \beta] &= \mathbb{E}[\mathbb{E}[\Delta_J \beta | \lambda]] \leq O_\eta(1) \cdot \mathbb{E} \left[ \mathbb{1}\{\Delta_J > 0\} \sum_{i=1}^d \min \{ \Delta_J^{-1}, \sqrt{\lambda_i} \} \right] \\ &\leq O_\eta(1) \cdot \mathbb{E} \left[ \sum_{i:\lambda_i = \infty, \Delta_i = 0} \frac{\mathbb{1}\{\Delta_J > 0\}}{\Delta_{\min}} + \sum_{i:\Delta_i > 0} \sqrt{\tau_i} \right] \leq O_\eta(1) \cdot \mathbb{E} \left[ \sum_{i:\Delta_i > 0} \sqrt{\tau_i} \right], \quad (8) \end{aligned}$$

where the last line follows since  $\mathbb{E}[J] = O_\eta(1)$  and

$$\mathbb{E} \left[ \sum_{i:\lambda_i = \infty, \Delta_i = 0} \frac{\mathbb{1}\{\Delta_J > 0\}}{\Delta_{\min}} \right] \leq \mathbb{E} \left[ \frac{J}{\Delta_{\min}} \right] \leq O_\eta(1) \cdot \frac{1}{\Delta_{\min}} \leq O_\eta(1) \max \{ i : \sqrt{\mathbb{E}[\tau_i]} \}.$$

The resulting is completed substituting  $\mathbb{E}[\sqrt{\tau_i}] \leq \sqrt{\mathbb{E}[\tau_i]}$  into Eq. (8) and applying Lemma 8 to show that  $\mathbb{E}[\tau_i] \leq O_\eta(1) \cdot \left(1 + \frac{\log(b_i)}{\Delta_i^2}\right)$ .  $\square$

#### Step 4: Putting it together

By substituting the bounds given in Lemma 10 into Eq. (7) and applying Lemma 8 we obtain

$$\begin{aligned} R_\mu^{\text{OCUCB-}n}(n) &\leq \sum_{i:\Delta_i>0} \Delta_i \mathbb{E}[\tau_i] + O_\eta(1) \cdot \sum_{i:\Delta_i>0} \sqrt{1 + \frac{\log(b_i)}{\Delta_i^2}} \\ &\leq O_\eta(1) \cdot \sum_{i:\Delta_i>0} \left( \Delta_i + \frac{1}{\Delta_i} \log \max \left\{ \frac{n\Delta_i^2 \log(n)}{k_{i,\rho}}, \log(n), e \right\} \right), \end{aligned}$$

which completes the proof of the finite-time bound.

**Asymptotic analysis.** Lemma 5 makes this straightforward. Let  $\varepsilon_n = \min\{\frac{\Delta_{\min}}{2}, \log^{-\frac{1}{4}}(n)\}$  and

$$\alpha_n = \min \left\{ \alpha : \inf_s \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log\left(\frac{\alpha}{Ks}\right)} \geq \mu_1 - \varepsilon_n \right\}.$$

Then by Lemma 5a with  $\rho = 1$  and  $\lambda_1, \dots, \lambda_K = \infty$  we have  $\sup_n \mathbb{E}[\alpha_n] \varepsilon_n^2 / K \leq O_\eta(1)$ . Then we modify the definition of  $\tau$  by

$$\tau_{i,n} = \min \left\{ s : \sup_{s' \geq s} \hat{\mu}_{i,s'} + \sqrt{\frac{2\eta}{s} \log(n \log(n))} \leq \mu_1 - \varepsilon_n \right\},$$

which is chosen such that if  $T_i(t-1) \geq \tau_{i,n}$ , then  $\gamma_i(t) \leq \mu_1 - \varepsilon_n$ . Therefore

$$R_\mu^{\text{OCUCB-}n}(n) \leq \Delta_{\max} \mathbb{E}[\alpha_n] + \sum_{i:\Delta_i>0} \Delta_i \mathbb{E}[\tau_{i,n}] \leq O_\eta(1) \cdot \frac{\Delta_{\max} K}{\varepsilon_n^2} + \sum_{i:\Delta_i>0} \Delta_i \mathbb{E}[\tau_{i,n}].$$

Classical analysis (eg., by Garivier and Cappé [2011]) shows that  $\limsup_{n \rightarrow \infty} \mathbb{E}[\tau_{i,n}] / \log(n) \leq 2\eta \Delta_i^{-2}$  and  $\lim_{n \rightarrow \infty} \varepsilon_n^{-2} / \log(n) = 0$ , which implies the asymptotic claim in Theorem 1.

$$\limsup_{n \rightarrow \infty} \frac{R_\mu^{\text{OCUCB-}n}(n)}{\log(n)} \leq \sum_{i:\Delta_i>0} \frac{2\eta}{\Delta_i}.$$

This naive calculation demonstrates a weakness of asymptotic results. The  $\Delta_{\max} K \varepsilon_n^{-2}$  term in the regret will typically dominate the higher-order terms except when  $n$  is outrageously large. A more careful argument (similar to the derivation of the finite-time bound) would lead to the same asymptotic bound via a nicer finite-time bound, but the details are omitted for readability. Interestingly the result is not dependent on  $\rho$  and so applies also to the MOSS-type algorithm that is recovered by choosing  $\rho = 0$ .

## 5 Discussion

The UCB family has a new member. This one is tuned for subgaussian noise and roughly mimics the OCUCB algorithm, but without needing advance knowledge of the horizon. The introduction of  $k_{i,\rho}$  is a minor refinement on previous measures of difficulty, with the main advantage being that it is very intuitive. The resulting algorithm is efficient and close to optimal theoretically. Of course there are open questions, some of which are detailed below.

**Shrinking the confidence level.** Empirically the algorithm improves significantly when the logarithmic terms in the definition of  $B_i(t-1)$  are dropped. There are several arguments that theoretically justify this decision. First of all if  $\rho > 1/2$ , then it is possible to replace the  $t \log(t)$  term in the definition of  $B_i(t-1)$  with just  $t$  and use part (a) of Lemma 5 instead of part (b). The price

is that the regret guarantee explodes as  $\rho$  tends to  $1/2$  (also not observed in practice). The second improvement is to replace  $\log(t)$  in the definition of  $B_i(t-1)$  with

$$\max \left\{ 0, \log \left( t \cdot \left( \sum_{j=1}^K \min \{ T_j(t-1), T_j(t-1)^\rho T_j(t-1)^{1-\rho} \} \right)^{-1} \right) \right\}, \quad (9)$$

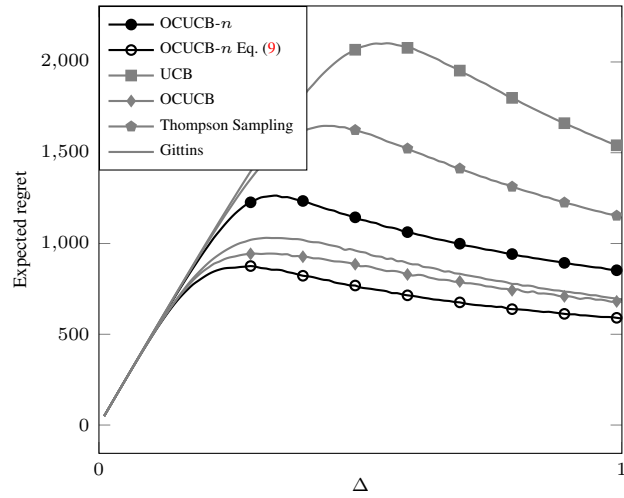
which boosts empirical performance and rough sketches suggest minimax optimality is achieved. I leave details for a longer article.

**Improving analysis and constants.** Despite its simplicity relative to OCUCB, the current analysis is still significantly more involved than for other variants of UCB. A cleaner proof would obviously be desirable. In an ideal world we could choose  $\eta = 1$  or (slightly worse) allow it to converge to 1 as  $t$  grows, which is the technique used in the KL-UCB algorithm [Cappé et al., 2013, and others]. I anticipate this would lead to an asymptotically optimal algorithm.

**Informational confidence bounds.** Speaking of KL-UCB, if the noise model is known more precisely (for example, it is bounded), then it is beneficial to use confidence bounds based on the KL divergence. Such bounds are available and could be substituted directly to improve performance without loss [Garivier, 2013, and others]. Repeating the above analysis, but exploiting the benefits of tighter confidence intervals would be an interesting (non-trivial) problem due to the need to exploit the non-symmetric KL divergences. It is worth remarking that confidence bounds based on the KL divergence are also *not* tight. For example, for Gaussian random variables they lead to the right exponential rate, but with the wrong leading factor, which in practice can improve performance as evidenced by the confidence bounds used by (near) Bayesian algorithms that exactly exploit the noise model (eg., Kaufmann et al. [2012], Lattimore [2016], Kaufmann [2016]). This is related to the “missing factor” in Hoeffding’s bound studied by Talagrand [1995].

**Precise lower bounds.** Perhaps the most important remaining problem for the subgaussian noise model is the question of lower bounds. Besides the asymptotic results by Lai and Robbins [1985] and Burnetas and Katehakis [1997] and the finite-time analysis in the two-armed case by Kulkarni and Lugosi [2000], there has been some recent progress on finite-time lower bounds, both in the OCUCB paper and a recent article by Garivier et al. [2016]. Results formalising the intuition in the introduction are given in Appendix A, but still there are regimes where the bounds are not matching.

**Empirical evaluation.** I plot the regret in the worst-case regime where  $K = 100$  and  $n = 5000$  and all sub-optimal arms have  $\Delta_i = \Delta > 0$ . Error bars are too small to see and all code/data will be made available with any final version. Each data point is the average of 14113 i.i.d. samples. There are two versions of OCUCB- $n$ , both with  $\eta = 1$ . The first using  $B_i$  as given after Eq. (1) and the second using Eq. (9). The performance is compared to UCB (as defined by Katehakis and Robbins [1995]), OCUCB ( $\alpha = 3, \psi = 2$ ), Thompson sampling and the finite-horizon Gittins index strategy. The two Bayesian algorithms use a flat Gaussian prior. Besides OCUCB- $n$ , only UCB and Thompson sampling do not depend on the horizon. A more extensive empirical evaluation would be quite interesting, especially when the gaps are unbalanced or even better with real data. Other synthetic evaluations were consistent with the results above, with OCUCB- $n$  performing comparably with the state-of-the-art in all regimes, and sometimes better. These results are omitted due to space constraints.





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## A Lower Bounds

I now prove a kind of lower bound showing that the form of the regret in Theorem 1 is approximately correct for  $\rho$  close to  $1/2$ . The result contains a lower order  $-\log \log(n)$  term, which for large  $n$  dominates the improvements, but is meaningful in many regimes.

**Theorem 11.** *Assume a standard Gaussian noise model and let  $\pi$  be any strategy and  $\mu \in [0, 1]^K$  be such that  $\frac{n\Delta_i^2}{k_{i,1/2} \log(n)} \geq 1$  for all  $i$ . Then one of the following holds:*

$$1. R_{\mu}^{\pi}(n) \geq \frac{1}{4} \sum_{i:\Delta_i>0} \frac{1}{\Delta_i} \log \left( \frac{n\Delta_i^2}{2k_{i,1/2} \log(n)} \right).$$

2. There exists an  $i$  with  $\Delta_i > 0$  such that

$$R_{\mu'}^{\pi}(n) \geq \frac{1}{2} \sum_{i:\Delta'_i>0} \frac{1}{\Delta'_i} \log \left( \frac{n\Delta_i'^2}{2k'_{i,1/2} \log(n)} \right)$$

where  $\mu'_i = \mu_i + 2\Delta_i$  and  $\mu'_j = \mu_j$  for  $j \neq i$  and  $\Delta'_i$  and  $k'_{i,\rho}$  are defined as  $\Delta_i$  and  $k_{i,\rho}$  but using  $\mu'$ .

*Proof.* On our way to a contradiction, assume that neither of the items hold. Let  $i$  be a suboptimal arm and  $\mu'$  be as in the second item above. I write  $\mathbb{P}'$  and  $\mathbb{E}'$  for expectation when when rewards are sampled from  $\mu'$ . Suppose

$$\mathbb{E}[T_i(n)] \leq \frac{1}{4\Delta_i^2} \log \left( \frac{n\Delta_i^2}{2k_{i,1/2} \log(n)} \right). \quad (10)$$

Then Lemma 2.6 in the book by [Tsybakov \[2008\]](#) and the same argument as used by [Lattimore \[2015\]](#) gives

$$\mathbb{P}\{T_i(n) \geq n/2\} + \mathbb{P}'\{T_i(n) < n/2\} \geq \frac{k_{i,1/2} \log(n)}{n\Delta_i^2} \equiv 2\delta.$$

By Markov's inequality

$$\mathbb{P}\{T_i(n) \geq n/2\} \leq \frac{2\mathbb{E}[T_i(n)]}{n} \leq \frac{1}{2n\Delta_i^2} \log \left( \frac{n\Delta_i^2}{2k_{i,1/2} \log(n)} \right) \leq \frac{\log(n)}{2n\Delta_i^2} \leq \delta.$$

Therefore  $\mathbb{P}'\{T_i(n) < n/2\} \geq \delta$ , which implies that

$$R_{\mu'}^{\pi}(n) \geq \frac{\delta n \Delta_i}{2} = \frac{1}{2} \sum_{j=1}^K \min \left\{ \frac{1}{\Delta_i}, \frac{1}{\Delta_j} \right\} \log(n) \geq \frac{1}{2} \sum_{j:\Delta'_j>0} \frac{1}{\Delta'_j} \log \left( \frac{n\Delta'_j}{2k'_{j,1/2} \log(n)} \right),$$

which is a contradiction. Therefore Eq. (10) does not hold for all  $i$  with  $\Delta_i > 0$ , but this also leads immediately to a contradiction, since then

$$R_{\mu}^{\pi}(n) = \sum_{i:\Delta_i>0} \Delta_i \mathbb{E}[T_i(n)] \geq \frac{1}{4} \sum_{i:\Delta_i>0} \frac{1}{\Delta_i} \log \left( \frac{n\Delta_i^2}{2k_{i,1/2} \log(n)} \right). \quad \square$$

## B Proof of Lemma 3

Monotonicity is obvious. Let  $\varepsilon > 0$  be such that  $\eta = 1 + 2\varepsilon$  and  $G_k = [(1 + \varepsilon)^k, (1 + \varepsilon)^{k+1}]$  and  $F_k = \left\{ \exists n \in G_k : S_n > \sqrt{2\eta n \log \max\{e, \log n\}} \right\}$ . Then

$$\mathbb{P}\left\{ \forall n : S_n \leq \sqrt{2\eta n \log \max\{e, \log n\}} \right\} = \mathbb{P}\left\{ \forall k \geq 0 : \neg F_k \right\} = \prod_{k=0}^{\infty} \mathbb{P}\left\{ \neg F_k \mid \neg F_1, \dots, \neg F_{k-1} \right\}.$$

Now we analyse the failure event  $F_k$ .

$$\begin{aligned} \mathbb{P}\{F_k \mid \neg F_1, \dots, \neg F_{k-1}\} &\leq \mathbb{P}\{F_k\} \\ &= \mathbb{P}\left\{ \exists n \in G_k : S_n > \sqrt{2\eta n \log \max\{e, \log n\}} \right\} \\ &\leq \exp \left( -\frac{2\eta(1 + \varepsilon)^k \log_+ \log(1 + \varepsilon)^k}{2(1 + \varepsilon)^{k+1}} \right) \\ &= \left( \frac{1}{k \log(1 + \varepsilon)} \right)^{1 + \frac{\varepsilon}{1 + \varepsilon}}. \end{aligned}$$

Since this is vacuous when  $k$  is small we need also need a naive bound.

$$\mathbb{P} \left\{ \exists n \in G_k : S_n \geq \sqrt{2\eta n \log \max \{e, \log n\}} \right\} \leq \exp(-\eta) < 1.$$

Combining these completes the results since for sufficiently large  $k_0$  (depending only on  $\eta$ ) we have that

$$p(\eta) \geq \exp(-\eta k_0) \prod_{k=k_0}^{\infty} (1 - \mathbb{P}\{F_k\}) \geq \exp(-\eta k_0) \prod_{k=k_0}^{\infty} \left( 1 - \left( \frac{1}{k \log(1+\varepsilon)} \right)^{1+\frac{\varepsilon}{1+\varepsilon}} \right) > 0.$$

### C Proof of Lemma 4

Let  $\alpha \geq 1$  be fixed and  $t_0 = \lceil 8\eta \log_+(b)/\Delta^2 \rceil$  and  $t_k = t_0 2^k$ . Then

$$\begin{aligned} \mathbb{P}\{\tau \geq \alpha t_0\} &\leq \mathbb{P}\{\exists t \geq \alpha t_0 : \hat{\mu}_t \geq \Delta/2\} \leq \sum_{k=0}^{\infty} \mathbb{P}\{\exists t \leq t_k : S_t \geq \alpha 2^{k-1} t_0 \Delta/2\} \\ &\leq \sum_{k=0}^{\infty} \exp\left(-\frac{\alpha^2 2^{2k-2} t_0^2 \Delta^2}{8\alpha 2^k t_0}\right) \leq \sum_{k=0}^{\infty} \exp\left(-\frac{\alpha 2^k}{4}\right) = O(\exp(-\alpha/4)). \end{aligned}$$

Therefore  $\mathbb{E}[(\tau/t_0)^2] = O(1)$  and so the result follows.

### D Proof of Lemma 5

Let  $\eta_1 = (1 + \eta)/2$  and  $\eta_2 = \eta/\eta_1$  and

$$\Lambda = \sum_{i=1}^d \min \left\{ \frac{1}{\Delta^2}, \frac{\lambda_i^\rho}{\Delta^{2-2\rho}} \log_+ \left( \frac{1}{\lambda_i \Delta^2} \right) \right\}.$$

Let  $x > 0$  be fixed and let  $G_k = [\eta_1^k, \eta_1^{k+1}]$ . We will use the peeling trick. First, by Lemma 2.

$$\begin{aligned} q_k &= \mathbb{P} \left\{ \inf_{s \in G_k} \hat{\mu}_s + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{x\Lambda}{\sum_{i=1}^d \min \{s^\rho, \lambda_i^\rho s^{1-\rho}\}} \right\}} \leq -\Delta \right\} \\ &\leq \mathbb{P} \left\{ \exists s \leq \eta_1^{k+1} : S_s + \sqrt{2\eta \eta_1^k \log \max \left\{ 1, \frac{x\Lambda}{\sum_{i=1}^d \min \{\eta_1^{k+1}, \lambda_i^\rho \eta_1^{(k+1)(1-\rho)}\}} \right\}} + \Delta \eta_1^k \leq 0 \right\} \\ &\stackrel{(a)}{\leq} \left( \frac{\sum_{i=1}^d \min \{\eta_1^{k+1}, \lambda_i^\rho \eta_1^{(k+1)(1-\rho)}\}}{x\Lambda} \right)^{\eta_2} \exp\left(-\frac{\Delta^2 \eta_1^{k-1}}{2}\right) \\ &= \left( \frac{\sum_{i=1}^d \min \{\eta_1^{k+1}, \lambda_i^\rho \eta_1^{(k+1)(1-\rho)}\}}{x\Lambda} \exp\left(-\frac{\Delta^2 \eta_1^k}{2\eta}\right) \right)^{\eta_2}, \end{aligned}$$

where (a) follows by Lemma 2. By the union bound

$$\begin{aligned} &\mathbb{P} \left\{ \inf_s \hat{\mu}_s + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{x\Lambda}{\sum_{i=1}^d \min \{s^\rho, \lambda_i^\rho s^{1-\rho}\}} \right\}} \leq -\Delta \right\} \leq \sum_{k=0}^{\infty} q_k \\ &\leq \sum_{k=0}^{\infty} \left( \frac{\sum_{i=1}^d \min \{\eta_1^{k+1}, \lambda_i^\rho \eta_1^{(k+1)(1-\rho)}\}}{x\Lambda} \exp\left(-\frac{\Delta^2 \eta_1^k}{2\eta}\right) \right)^{\eta_2} \\ &\leq \left( \frac{1}{x\Lambda} \sum_{i=1}^d \sum_{k=0}^{\infty} \min \{\eta_1^{k+1}, \lambda_i^\rho \eta_1^{(k+1)(1-\rho)}\} \exp\left(-\frac{\Delta^2 \eta_1^k}{2\eta}\right) \right)^{\eta_2} \\ &= O\left(\frac{\eta}{\eta-1}\right) \cdot x^{-\eta_2}, \end{aligned}$$

where the last line follows from Lemma 12. Therefore  $\mathbb{P}\{\alpha \geq x\Lambda\} \leq O\left(\frac{\eta}{\eta-1}\right) \cdot x^{-\eta_2}$ .

Now the first part follows easily since  $\mathbb{E}[\alpha] \leq \int_0^\infty \mathbb{P}\{\alpha \geq x\Lambda\} = O\left(\frac{\eta}{(\eta-1)^2}\right) \cdot \Lambda$ . Therefore

$$\begin{aligned} \Delta\mathbb{E}[\alpha] &\leq O\left(\frac{\eta}{(\eta-1)^2}\right) \cdot \sum_{i=1}^d \min\left\{\frac{1}{\Delta}, \lambda_i^\rho \Delta^{2\rho-1} \log_+\left(\frac{1}{\lambda_i \Delta^2}\right)\right\} \\ &\leq O\left(\frac{\eta}{(2\rho-1)(\eta-1)^2}\right) \cdot \sum_{i=1}^d \min\left\{\frac{1}{\Delta}, \sqrt{\lambda_i}\right\}. \end{aligned}$$

For the second part let  $x_0 = \Lambda / \text{productlog}(\Lambda)$  where productlog is the inverse of the function  $x \rightarrow x \exp(x)$ .

$$\begin{aligned} \mathbb{E}[\beta] &\leq \int_0^\infty \mathbb{P}\{\beta \geq x\} dx \leq x_0 + \int_{x_0}^\infty \mathbb{P}\left\{\alpha \geq \frac{x}{\Lambda} \log(x)\right\} dx \\ &\leq x_0 + O\left(\frac{\eta}{\eta-1}\right) \cdot \int_{x_0}^\infty \left(\frac{\Lambda}{x \log(x)}\right)^{\eta_2} dx \\ &\leq x_0 + O\left(\frac{\eta}{\eta-1}\right) \cdot \left(\frac{\Lambda}{\log(x_0)}\right)^{\eta_2} \int_{x_0}^\infty x^{-\eta_2} dx \\ &\leq x_0 + O\left(\frac{\eta}{(\eta-1)^2}\right) \cdot \left(\frac{\Lambda}{\log(x_0)}\right)^{\eta_2} x_0^{1-\eta_2} = O\left(\frac{\eta}{(\eta-1)^2}\right) \cdot \frac{\Lambda}{\text{productlog}(\Lambda)}. \end{aligned}$$

If  $\Lambda < e$ , then the result is trivial. For  $\Lambda \geq e$  we have  $\text{productlog}(\Lambda) \geq 1$ . Then

$$\begin{aligned} \Delta\mathbb{E}[\beta] &\leq O\left(\frac{1}{(\eta-1)^2}\right) \cdot \frac{\Delta\Lambda}{\text{productlog}(\Lambda)} \\ &\leq O\left(\frac{1}{(\eta-1)^2}\right) \cdot \sum_{i=1}^d \min\left\{\frac{1}{\Delta}, \frac{\lambda_i^\rho \Delta^{2\rho-1}}{\text{productlog}(\Lambda)} \log_+\left(\frac{1}{\lambda_i \Delta^2}\right)\right\}. \end{aligned}$$

By examining the inner minimum we see that if  $\Delta \geq \lambda_i^{-\frac{1}{2}}$ , then  $1/\Delta \leq \lambda_i^{\frac{1}{2}}$ . If  $\Delta < \lambda_i^{-\frac{1}{2}}$ , then

$$\begin{aligned} \min\left\{\frac{1}{\Delta}, \frac{\lambda_i^\rho \Delta^{2\rho-1}}{\text{productlog}(\Lambda)} \log_+\left(\frac{1}{\lambda_i \Delta^2}\right)\right\} &< \frac{\lambda_i^{\frac{1}{2}}}{\max\{1, \text{productlog}(\Delta^{-2})\}} \log_+\left(\frac{1}{\lambda_i \Delta^2}\right) \\ &\leq 2\lambda_i^{\frac{1}{2}}. \end{aligned}$$

Therefore  $\mathbb{E}[\Delta t] \leq O\left(\frac{\eta}{(\eta-1)^2}\right) \cdot \sum_{i=1}^d \min\{\Delta^{-1}, \sqrt{\lambda_i}\}$  as required.

## E Proof of Lemma 8

Since  $J$  is sub-exponentially distributed with rate dependent only on  $\eta$  we have  $\sqrt{\mathbb{E}[J^2]} = O(1)$ . By using Lemma 4 we obtain

$$\begin{aligned} \sqrt{\mathbb{E}[\tau_i^2]} &= \sqrt{\mathbb{E}[\mathbb{E}[\tau_i^2 | K_{i,\rho}]]} \\ &= O_\eta(1) \cdot \sqrt{\mathbb{E}\left[\left(1 + \frac{1}{\Delta_i^2} \log(B_i)\right)^2\right]} = O_\eta(1) \cdot \left(1 + \frac{1}{\Delta_i^2} \log(b_i)\right). \end{aligned}$$

The latter inequality follows by noting that  $B_i \geq e$  and  $(1 + c \log(x))^2$  is concave for  $x \geq e$  and  $c > 0$ .

$$\begin{aligned} \sqrt{\mathbb{E} \left[ \left( 1 + \frac{1}{\Delta_i^2} \log(B_i) \right)^2 \right]} &\leq 1 + \frac{1}{\Delta_i^2} \log(\mathbb{E}[B_i]) \\ &= 1 + \frac{1}{\Delta_i^2} \log \left( \mathbb{E} \left[ \max \left\{ \log(n), \frac{n\Delta_i^2 \log(n)}{K_{i,\rho}} \right\} \right] \right) \\ &= O_\eta(1) \cdot \left( 1 + \frac{1}{\Delta_i^2} \log \left( \max \left\{ \log(n), \frac{n\Delta_i^2 \log(n)}{k_{i,\rho}} \right\} \right) \right), \end{aligned}$$

where the last inequality follows from (a)  $K_{i,\rho} \geq 1$  and (b) Azuma's concentration inequality implies that  $\mathbb{P}\{K_{i,\rho} \leq c_\eta \rho(\eta) k_{i,\rho}/2\} = O(k_{i,\rho}^{-1})$  as shown in the following appendix. Finally by Holder's inequality

$$\mathbb{E}[J\tau_i] \leq \sqrt{\mathbb{E}[J^2]\mathbb{E}[\tau_i^2]} \leq O_\eta(1) \cdot \left( 1 + \frac{1}{\Delta_i^2} \log(b_i) \right).$$

## F Tail Bound on $K_{i,\rho}$

Recall that  $K_{i,\rho} = 1 + c_\eta \sum_{j \in \Phi, j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\}$  and  $k_{i,\rho} = 1 + \sum_{j \neq i}^K \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\}$ . Therefore by Azuma's inequality and naive simplification we have

$$\begin{aligned} \mathbb{P}\{K_{i,\rho} \leq c_\eta \rho(\eta) k_{i,\rho}/2\} &\leq \mathbb{P} \left\{ \sum_{j \in \Phi, j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\} \leq \frac{\rho(\eta)}{2} \sum_{j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\} \right\} \\ &\stackrel{(a)}{\leq} \exp \left( - \frac{\left( \rho(\eta) \sum_{j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\} \right)^2}{2 \sum_{j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\}^2} \right) \\ &\stackrel{(b)}{\leq} \exp \left( - \frac{\rho(\eta)^2 \sum_{j \neq i} \min \left\{ 1, \Delta_i^{2\rho} / \Delta_j^{2\rho} \right\}}{2} \right) \\ &\stackrel{(c)}{=} O(k_{i,\rho}^{-1}), \end{aligned}$$

where (a) follows from Azuma's inequality and (b) since  $\min\{1, x\}^2 \leq \min\{1, x\}$  and (c) by  $\exp(-x) \leq 1/x$  for all  $x \geq 0$ .

## G Proof of Lemma 10b

Recall that we are trying to show that

$$\mathbb{E} \left[ \sum_{i: \Delta_i < \Delta_J/4} \Delta_i \tau_{\Delta_i} \right] = O \left( \sum_{i: \Delta_i > 0} \Delta_i \mathbb{E}[J\tau_i] \right). \quad (11)$$

Let  $E$  be the event that  $\Delta_2 \leq \Delta_J/4$  and define random sets  $A_1 = \{i : \Delta_i \in (2\Delta_J, \infty)\}$  and  $A_2 = \{i : \Delta_i \in [\Delta_J, 2\Delta_J]\}$ . For  $i \in A_1$  we have  $\Delta_i > 2\Delta_J$  and since  $J \in \Phi$  we have  $\gamma_{J,J}(t) \geq \mu_J \geq \mu_1 - \Delta_i/2$ . Therefore  $i \in A_1$  implies that  $\tau_{\Delta_i} = 1$  and so  $T_i(n) \leq \tau_i$ . Let  $\lambda \in (0, \infty]^K$  be given by  $\lambda_i = \tau_i$  for  $i \in A_1$  and  $\lambda_i = \infty$  otherwise.

$$\alpha = \min \left\{ \alpha : \inf_s \hat{\mu}_{2,s} + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{\alpha}{\sum_{i=1}^K \min \{s, \lambda_i^\rho s^{1-\rho}\}} \right\}} \geq \mu_2 - \frac{\Delta_J}{4} \right\}.$$

It is important to note that we have used  $\hat{\mu}_{2,s}$  in the definition of  $\alpha$  and not  $\hat{\mu}_{1,s}$  that appeared in the proof of part (a) of this lemma. The reason is to preserve independence when samples from the

first arm are used later. Let  $\beta = \min \{\beta \geq 0 : \beta \log(\beta) = \alpha\}$ . If  $E$  holds, then for  $t \geq \beta$  we have  $\gamma_2(t) \geq \mu_2 - \Delta_J/4 \geq \mu_1 - \Delta_J/2$ , which implies that

$$\mathbb{1}\{E\} \sum_{i \in A_2} T_i(n) \leq t_{\Delta_J} + \sum_{i \in A_2} \tau_i \leq \beta + \sum_{i \in A_2} \tau_i.$$

Therefore for any  $s, t \leq n$  the concavity of  $\min \{s, \cdot\}$  and  $x \rightarrow x^\rho$  combined with Jensen's inequality implies that

$$\mathbb{1}\{E\} \sum_{i \in A_2} \min \{s, T_i(t-1)^\rho s^{1-\rho}\} \leq \sum_{i \in A_2} \min \left\{ s, \left( \frac{\beta + \sum_{i \in A_2} \tau_i}{|A_2|} \right)^\rho s^{1-\rho} \right\}.$$

We are getting close to an application of Lemma 5. Let  $\omega \in (0, \infty]^K$  be given by

$$\omega_j = \begin{cases} \tau_j & \text{if } j \in A_1 \\ \beta/|A_2| + \sum_{j \in A_2} \tau_j/|A_2| & \text{if } j \in A_2 \\ \infty & \text{otherwise,} \end{cases}$$

which has been chosen such that for  $T_1(t-1) = s$  and if  $E$  holds, then

$$\begin{aligned} B_1(t-1) &\geq \max \left\{ 1, \frac{t \log(t)}{\sum_{j=1}^K \min \{s, T_j(t-1)^\rho s^{1-\rho}\}} \right\} \\ &\geq \max \left\{ 1, \frac{t \log(t)}{\sum_{j=1}^K \min \{s, \omega_j^\rho s^{1-\rho}\}} \right\}. \end{aligned} \quad (12)$$

Now let  $i$  be the index of some arm for which  $\Delta_i < \Delta_J/4$  and define

$$\alpha_i = \min \left\{ \alpha : \inf_s \hat{\mu}_{1,s} + \sqrt{\frac{2\eta}{s} \log \max \left\{ 1, \frac{\alpha}{\sum_{j=1}^K \min \{s, \omega_j^\rho s^{1-\rho}\}} \right\}} \geq \mu_1 - \frac{\Delta_i}{2} \right\}$$

and  $\beta_i = \min \{\beta \geq 0 : \beta \log(\beta) = \alpha_i\}$ . Therefore by Eq. (12), if  $E$  holds and  $t \geq \beta_i$ , then  $\gamma_1(t) \geq \mu_1 - \Delta_i/2$  and so  $t_{\Delta_i} \leq \beta_i$ . At last we are able to write  $t_{\Delta_i}$  in terms of something for which the expectation can be controlled.

$$\begin{aligned} \mathbb{E} \left[ \sum_{i: \Delta_i < \Delta_J/4} \Delta_i \tau_{\Delta_i} \right] &\leq \mathbb{E} \left[ \sum_{i: \Delta_i < \Delta_J/4} \Delta_i \beta_i \right] \\ &\leq O_\eta(1) \cdot \mathbb{E} \left[ \sum_{i: \Delta_i < \Delta_J/4} \sum_{j=1}^K \min \left\{ \frac{1}{\Delta_i}, \sqrt{\omega_j} \right\} \right] \\ &\leq O_\eta(1) \cdot \mathbb{E} \left[ \sum_{i: \Delta_i < \Delta_J/4} \left( \sum_{j \in A_1} \sqrt{\tau_j} + |A_2| \sqrt{\omega_j} + \frac{J}{\Delta_{\min}} \right) \right] \\ &\leq O_\eta(1) \cdot \mathbb{E} \left[ \sum_{j \in A_1} J \sqrt{\tau_j} + \frac{J^2}{\Delta_{\min}} + J |A_2| \sqrt{\omega_j} \right]. \end{aligned} \quad (13)$$

The first two terms are easily bounded as we shall soon see. For the last we have

$$\begin{aligned} \mathbb{E} [J |A_2| \sqrt{\omega_j}] &\leq O_\eta(1) \cdot \sqrt{\mathbb{E} [|A_2|^2 \omega_j]} = O_\eta(1) \cdot \sqrt{\mathbb{E} \left[ |A_2| \sum_{j \in A_2} \tau_j + |A_2| \beta \right]} \\ &\leq O_\eta(1) \cdot \left( \sqrt{\mathbb{E} \left[ |A_2| \sum_{j \in A_2} \tau_j \right]} + \sqrt{\mathbb{E} [|A_2| \beta]} \right) \end{aligned} \quad (14)$$

Bounding each term separately. For the first, let  $\tilde{A}_\ell = \{j : \Delta_j \in [2^\ell, 2^{\ell+2})\}$ , which is chosen such that no matter the value of  $\Delta_j$  there exists an  $\ell \in \mathbb{Z}$  with  $A_2 \subseteq A_\ell$ .

$$\begin{aligned} \sqrt{\mathbb{E} \left[ |A_2| \sum_{j \in A_2} \tau_j \right]} &\leq O(1) \cdot \sqrt{\sum_{\ell \in \mathbb{Z}} |\tilde{A}_\ell| \sum_{j \in \tilde{A}_\ell} \mathbb{E}[\tau_j]} \\ &\leq O(1) \cdot \sqrt{\sum_{\ell \in \mathbb{Z}} |\tilde{A}_\ell|^2 \max_{j \in \tilde{A}_\ell} \mathbb{E}[\tau_j]} \\ &\leq O_\eta(1) \cdot \sum_{j: \Delta_j > 0} \sqrt{1 + \frac{\log(b_j)}{\Delta_j^2}}, \end{aligned} \quad (15)$$

where the last inequality follows because  $\sum_{\ell \in \mathbb{Z}} \mathbb{1}\{j \in \tilde{A}_\ell\} = 2$  for each  $j$  and from Lemma 8, which gives the same order-bound on  $\mathbb{E}[\tau_j]$  for all  $j \in \tilde{A}_\ell$  for fixed  $\ell$ . For the second term in Eq. (14) we have

$$\begin{aligned} \mathbb{E} \left[ \sqrt{|A_2| \beta} \right] &\stackrel{(a)}{\leq} O_\eta(1) \cdot \mathbb{E} \left[ \sqrt{|A_2| \left( \sum_{j: \lambda_j = \infty} \frac{1}{\Delta_j^2} + \frac{1}{\Delta_J} \sum_{j: \lambda_j < \infty} \sqrt{\tau_j} \right)} \right] \\ &\stackrel{(b)}{\leq} O_\eta(1) \cdot \left( \sqrt{\mathbb{E} \left[ |A_2| \sum_{j: \lambda_j = \infty} \frac{1}{\Delta_j^2} \right]} + \mathbb{E} \left[ \frac{|A_2|}{\Delta_J} \right] + \sum_{j: \Delta_j > 0} \mathbb{E}[\sqrt{\tau_j}] \right) \\ &\stackrel{(c)}{\leq} O_\eta(1) \sum_{j: \Delta_j > 0} \sqrt{1 + \frac{\log(b_j)}{\Delta_j^2}}, \end{aligned}$$

where (a) follows from Lemma 5 and (b) since for all  $x, y \geq 0$  it holds that  $\sqrt{x+y} \leq \sqrt{x} + \sqrt{y}$  and  $\sqrt{xy} \leq x+y$ . To get (c) we bound the first term as in Eq. (15), the second by the fact that arms in  $j \in A_2$  have  $\Delta_j \leq 2\Delta_J$  and the third using Lemma 8. Finally by substituting this into Eq. (13) we have

$$\begin{aligned} \mathbb{E} \left[ \sum_{i \in A_3} \Delta_i \tau_{\Delta_i} \right] &\leq O_\eta(1) \cdot \left( \mathbb{E} \left[ \sum_{j \in A_1} J \sqrt{\tau_j} + \frac{J^2}{\Delta_{\min}} \right] + \sum_{j: \Delta_j > 0} \left( 1 + \frac{\log(b_j)}{\Delta_j^2} \right) \right) \\ &\leq O_\eta(1) \sum_{j: \Delta_j > 0} \left( 1 + \frac{\log(b_j)}{\Delta_j^2} \right), \end{aligned}$$

where the last line follows since  $\mathbb{E}[J^2/\Delta_{\min}] = O_\eta(1)\Delta_{\min}^{-1}$  and by Lemma 8

$$\mathbb{E}[J\sqrt{\tau_j}] \leq \sqrt{\mathbb{E}[J^2]\mathbb{E}[\tau_j]} = O_\eta(1) \cdot \sqrt{1 + \frac{\log(b_j)}{\Delta_j^2}},$$

which completes the proof.

## H Technical Lemmas

**Lemma 12.** *Let  $\eta > 1$  and  $\rho \in [0, 1]$  and  $\lambda \in (0, \infty]$  and  $x > 0$ , then*

$$\begin{aligned} \sum_{k=0}^{\infty} \min \left\{ \eta^{k+1}, \lambda^\rho \eta^{(1-\rho)(k+1)} \right\} \exp(-x\eta^k) &\leq \begin{cases} \frac{1}{x} \left( \frac{2}{e} + \frac{\eta}{\log(\eta)} \right) & \text{if } x\lambda \geq 1 \\ \frac{\lambda^\rho x^{\rho-1}}{\log(\eta)} \left( 1 + \frac{1}{e} + \log\left(\frac{1}{\lambda x}\right) \right) & \text{otherwise.} \end{cases} \\ &= O \left( \frac{\eta}{\eta-1} \right) \cdot \min \left\{ \frac{1}{x}, \lambda^\rho x^{\rho-1} \log_+ \left( \frac{1}{\lambda x} \right) \right\}. \end{aligned}$$

*Proof.* Let  $f(k) = \min \{ \eta^{k+1}, \lambda^\rho \eta^{(k+1)(1-\rho)} \} \exp(-x\eta^k)$ , which is unimodal and so  $\sum_{k=0}^{\infty} f(k) \leq 2 \sup_k f(k) + \int_0^{\infty} f(k) dk$ . If  $x\lambda \geq 1$ , then

$$\int_0^{\infty} f(k) dk \leq \eta \int_0^{\infty} \eta^k \exp(-\eta^k x) dk = \frac{\eta}{x \log(\eta)}.$$

If  $x\lambda < 1$ , then let  $k_\lambda$  be such that  $\eta^{k_\lambda} = \lambda^\rho \eta^{k_\lambda(1-\rho)}$  and  $k_x$  be such that  $\eta^{k_x} = 1/x$ .

$$\begin{aligned} \int_0^{\infty} f(k) dk &\leq \eta \int_0^{k_\lambda} \eta^k dk + \eta \int_{k_\lambda}^{k_x} \lambda^\rho \eta^{k_x(1-\rho)} dk + \eta \int_{k_x}^{\infty} \lambda^\rho \eta^{k(1-\rho)} \exp(-x\eta^k) dk \\ &= \frac{\lambda-1}{\log(\eta)} + \eta(k_x - k_\lambda) \lambda^\rho x^{\rho-1} + \eta \lambda^\rho x^{\rho-1} \int_{k_x}^{\infty} \eta^{(k-k_x)(1-\rho)} \exp(-\eta^{k_x-k}) dk \\ &\leq \frac{\lambda-1}{\log(\eta)} + \frac{\eta \lambda^\rho x^{\rho-1} \log(\frac{1}{\lambda x})}{\log(\eta)} + \frac{\eta \lambda^\rho x^{\rho-1}}{e \log(\eta)} \end{aligned}$$

Finally

$$\sup_k f(k) \leq \min \left\{ \frac{1}{ex}, \eta \lambda^\rho x^{\rho-1} \right\}.$$

Therefore

$$\sum_{k=0}^{\infty} \min \left\{ \eta^{k+1}, \lambda^\rho \eta^{(1-\rho)(k+1)} \right\} \exp(-x\eta^k) \leq \begin{cases} \frac{1}{x} \left( \frac{2}{e} + \frac{\eta}{\log(\eta)} \right) & \text{if } x\lambda \geq 1 \\ \frac{\lambda^\rho x^{\rho-1}}{\log(\eta)} \left( 1 + \frac{1}{e} + \log\left(\frac{1}{\lambda x}\right) \right) & \text{otherwise.} \end{cases}$$

□



## I Table of Notation

$K$	number of arms
$n$	horizon
$t$	current time step
$\xi_t$	noise in time step $t$
$\eta$	constant parameter greater than 1 determining width of confidence interval
$\rho$	constant parameter in $[1/2, 1]$
$\eta_1, \eta_2$	$\eta_1 = (1 + \eta)/2$ and $\eta_2 = \eta/\eta_1$
$\mu_i$	expected return of arm $i$
$\hat{\mu}_{i,s}$	empirical estimate of return of arm $i$ based on $s$ samples
$\hat{\mu}_i(t)$	empirical estimate of return of arm $i$ after time step $t$
$\Delta_i$	gap between the expected returns of the best arm and the $i$ th arm
$\Delta_{\min}$	minimal non-zero gap $\Delta_{\min} = \min \{\Delta_i : \Delta_i > 0\}$
$\Delta_{\max}$	maximum gap $\Delta_{\max} = \max_i \Delta_i$
$\log_+(x)$	$\max \{1, \log(x)\}$
$B_i$ and $b_i$	see Eq. (5)
$k_{i,\rho}$	$\sum_{j=1}^K \min\{1, \Delta_i^{2\rho} / \Delta_j^{2\rho}\}$
$K_{i,\rho}$	see Eq. (4)
$\tau_i$	see Eq. (5)
$\tau_\Delta$	see Eq. (6)
$p(\eta)$	see Lemma 3
$\Phi$	set of optimistic arms Eq. (3)
$J$	$J = \min \Phi$